



EEL3701

Menu

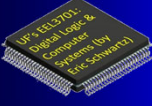
- Boolean Algebra
 - > Theorems
 - > Definitions




See examples on web:
[BooleanAlgebra.PDF](#)

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EEL3701 The Mathematics of Logic Design - Boolean Algebra

- **Basic Postulates & Theorems**
 - > **Identity Laws**

| | |
|----------------------|-----------------|
| 1. $\bar{X} + 0 = X$ | $X \cdot 1 = X$ |
| 2. $X + 1 = 1$ | $X \cdot 0 = 0$ |
 - > **Idempotent Laws**

| | |
|----------------|-----------------|
| 3. $X + X = X$ | $X \cdot X = X$ |
|----------------|-----------------|
 - > **Involution & Complementarity Laws**

| | |
|-----------------|------------------|
| 4. $(X')' = X$ | |
| 5. $X + X' = 1$ | $X \cdot X' = 0$ |
 - > **Commutative Laws**

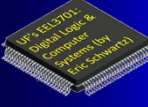
| | |
|--------------------|-------------------------|
| 6. $X + Y = Y + X$ | $X \cdot Y = Y \cdot X$ |
|--------------------|-------------------------|
 - > **Associative Laws**

| | |
|----------------------------|---------------------------------|
| 7. $(X+Y)+Z=X+(Y+Z)=X+Y+Z$ | $(XY) \cdot Z=X \cdot (YZ)=XYZ$ |
|----------------------------|---------------------------------|
 - > **Absorption Laws**

| | |
|--------------------------|-----------------------|
| 8. $X \cdot (X + Y) = X$ | $X + (X \cdot Y) = X$ |
|--------------------------|-----------------------|

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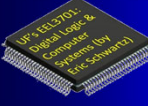
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- **Important Theorems**
 - > **Distributive Laws**
 9. $X \cdot (Y + Z) = X \cdot Y + X \cdot Z$ $X + (Y \cdot Z) = (X + Y) \cdot (X + Z)$
 - > **De Morgan's Law**
 10. $(X + Y + Z)' = X' \cdot Y' \cdot Z'$ $(X \cdot Y \cdot Z)' = X' + Y' + Z'$
 - > **Duality**
 If we treat **+ and •** as dual pairs, and **0 and 1** as dual pairs, the theorems in one column can be deduced from the other, e.g.,

| | | |
|---------------------|---------------------|----------------------|
| 1. $X + 0 = X$ | 2. $X + 1 = 1$ | 5. $X + X' = 1$ |
| ↓ ↓ ↓ ↓ | ↓ ↓ ↓ ↓ | ↓ ↓ ↓ ↓ |
| 1D. $X \cdot 1 = X$ | 2D. $X \cdot 0 = 0$ | 5D. $X \cdot X' = 0$ |

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EEL3701 Truth Table Proof of DeMorgan's Law (and use of Duality)

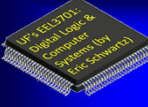
$\neg(A + B) = \neg A \cdot \neg B \xleftrightarrow{\text{Dual}} \neg(A \cdot B) = \neg A + \neg B$

| A | B | $\neg(A+B)$ | $\neg A \cdot \neg B$ | $\neg(A \cdot B)$ | $\neg A + \neg B$ |
|---|---|-------------|-----------------------|-------------------|-------------------|
| 0 | 0 | 1 | 1 | 1 | 1 |
| 0 | 1 | 0 | 0 | 1 | 1 |
| 1 | 0 | 0 | 0 | 1 | 1 |
| 1 | 1 | 0 | 0 | 0 | 0 |



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An Example of DeMorgan's Law

Problem: Simplify $\overline{[(X + YZ) + W]}$

Solutions: Use DeMorgan's Law over and over 

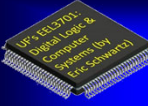
- Let $(X + YZ) = A$ to make variables easier to work with

1. Substituting: $\overline{[(X + YZ) + W]} = \overline{[A+W]}$
 > By DeMorgan's Law, $= \overline{[A+W]} = \overline{A} \cdot \overline{W}$
2. But $\overline{A} = \overline{(X + YZ)}$
 > By DeMorgan's Law, $= \overline{(X + YZ)} = \overline{X} \cdot \overline{(YZ)}$
 > By DeMorgan's Law, $= \overline{(YZ)} = \overline{Y} + \overline{Z}$
3. Putting 2 all together:
 > $\overline{A} = \overline{(X + YZ)} = \overline{X} \cdot \overline{(YZ)} = \overline{X} \cdot (\overline{Y} + \overline{Z})$
4. Substituting 3 into 1 \Rightarrow
 > Answer $= \overline{X} \cdot (\overline{Y} + \overline{Z}) \cdot \overline{W}$

$(B+C)' = B' \cdot C'$
 $(B \cdot C)' = B' + C'$

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An Example of DeMorgan's Law

Problem: Simplify $\overline{[(X + YZ) + W]}$

Solutions: Use DeMorgan's Law over and over

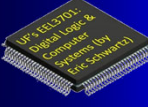
4. Substituting 3 into 1 \Rightarrow
 > Answer $= \overline{X} \cdot (\overline{Y} + \overline{Z}) \cdot \overline{W}$
5. By the *distributive law*,
 > Answer $= \overline{X} \cdot \overline{Y} \cdot \overline{W} + \overline{X} \cdot \overline{Z} \cdot \overline{W}$
6. Omitting dots for clarity,
 > Answer $= \overline{X} \overline{Y} \overline{W} + \overline{X} \overline{Z} \overline{W}$

Preferred Answer $= \overline{W} \overline{X} \overline{Y} + \overline{W} \overline{X} \overline{Z}$

> Above is in **Lexical Order**

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EEL3701 Boolean Expressions in Lexical Order

- Consider $XYZ=XZY=YXZ=YZX=ZXY=ZYX$ and $X+Y+Z=X+Z+Y=Y+X+Z=Y+Z+X=Z+X+Y=Z+Y+X$.
 > Since the terms can be written in many ways (by associativity and/or commutativity), we need conventions.

Rule H1: You shall write all the Boolean expressions in **lexical** order, /A, A, /B, B,....., /Z, Z. Then each term is in **word** order.

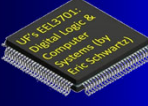
- The answer of the previous example should, therefore, be:

$$/W /X /Y + /W /X /Z$$
 rather than

$$/X /Y /W + /X /Z /W \text{ or } /Y /X /W + /Z /W /X$$
 (or many other possibilities)

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EEL3701 Boolean Expressions in Lexical Order

- Consider $/D_2 D_3 D_0 /D_1 = D_3 /D_2 /D_1 D_0 = D_0 /D_1 /D_2 D_3$
 > Since the terms can be written in many ways (by associativity and/or commutativity), we need conventions.

Rule H2: You shall write all the Boolean expressions with same name in **reverse subscript numerical** order. Especially important within a single product (AND) term.

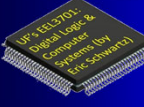
- For example, write the above product term as

$$D_3 /D_2 /D_1 D_0$$
- For example, write

$$E A_{13} /A_{14} + /A_{15} \text{ as } /A_{15} + /A_{14} A_{13} E$$

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EEL3701 SOP (Sum of Products)

Definitions: Any function can be written as a sum of products that is as a disjunction of conjunctive terms (OR of ANDs). A minterm is a conjunctive (AND) term that is 1 in only 1 row of a truth table.

| a | b | c | m ₇ | m ₆ | m ₅ | m ₄ | m ₃ | m ₂ | m ₁ | m ₀ | f(a,b,c) |
|---|---|---|----------------|----------------|----------------|----------------|----------------|----------------|----------------|----------------|----------|
| 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 1 |
| 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 |
| 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 1 |
| 0 | 1 | 1 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 1 |
| 1 | 0 | 0 | 0 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 |
| 1 | 0 | 1 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 1 |
| 1 | 1 | 0 | 0 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |
| 1 | 1 | 1 | 1 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 |

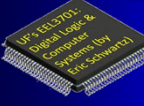
$$f = \sum_i \text{Minterm}_i$$

$f(a,b,c) = m_0 + m_2 + m_3 + m_5 + m_6$
 $= /a/b/c + /a b /c + /a b c + a /b c + a b /c$

Q: Can this be simplified? _____ Why? _____ How? _____

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EEL3701 Minimum Sum Of Product (MSOP)

- When you cannot reduce an SOP further, this is called a minimum sum of products (MSOP) expression

Ex: Simplify the following expression (from previous page)

$$f(a,b,c) = m_0 + m_2 + m_3 + m_5 + m_6 = /a/b/c + /a b/c + /a b c + a/b c + a b/c$$

We quite often use $X+/X = 1$, $X \cdot 1 = X$, $X+X = X$, & $X \cdot X = X$.

Thus,

$$m_0 + m_2 = /a/b/c + /a b/c = /a /c (b + /b) = /a /c$$

$$m_2 + m_3 = /a b/c + /a b c = /a b (/c + c) = /a b$$

$$m_2 + m_6 = /a b/c + a b/c = (/a + a) b/c = b/c$$

$$f(a,b,c) = m_0 + m_2 + m_3 + m_5 + m_6 = m_0 + m_2 + m_2 + m_2 + m_3 + m_5 + m_6$$

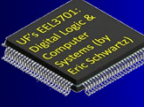
$$= (m_0 + m_2) + (m_2 + m_3) + (m_2 + m_6) + m_5$$

$$= /a /c + /a b + b/c + a/b c$$

$$f(a,b,c) = /a b + /a /c + a /b c + b /c$$

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EEL3701 Algebraic Simplification - Boolean Algebra

Simplification using the axioms of Boolean Algebra

EX: $G_{SOP} = \overline{A} \overline{B} + A \overline{B} = (\overline{A} + A) \overline{B} = \overline{B}$

$\overline{X} + X = 1$
 $X \cdot 1 = X$

Note: Truth table for $\overline{A} \overline{B} + A \overline{B}$ is same as for \overline{B}

Let's view similar product of sums (POS) result

EX: $G_{POS} = (\overline{A} + \overline{B})(A + \overline{B})$

Commutative Law

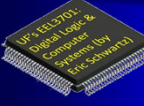
$= (\overline{B} + \overline{A})(\overline{B} + A)$

$G_{POS} = \overline{B}$

Apply Distributive Law: $(X+Y)(X+Z) = X+(YZ)$
 So $(X+Y)(X+Y) = X+(Y/Y) = X+0 = X$
 Now let $X=\overline{B}, Y=\overline{A} \Rightarrow \overline{Y} = A$

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EEL3701 Algebraic Simplification for Simplified Circuits

- We want to simplify (and eliminate) literals (terms) because each literal is a gate input. The smaller the number of literals the better!

EX: Any activation levels

$G_{SOP} = \overline{A} \overline{B} + A \overline{B}$

$G_{SOP} = \overline{B}$

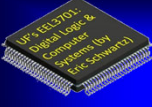
Saved 4 gates!

$G_{SOP} = \overline{B}$

Wires are cheap!

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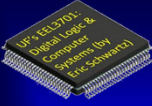
More Time?

- If time, demonstrate more
 - > Mixed-logic examples
 - > Signed binary numbers
 - > Boolean Arithmetic
 - > Boolean Algebra

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The End!

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